

Disclosure of Endogenous Information

Matthew Gentzkow and Emir Kamenica*
University of Chicago

August 2012

Abstract

We study the effect of disclosure requirements in environments where experts publicly acquire private information before engaging in a persuasion game with a decision-maker. In contrast to settings where private information is exogenous, we show that disclosure requirements do not change the set of pure-strategy equilibrium outcomes regardless of the players' preferences.

Keywords: persuasion; regulation; verifiable types; asymmetric information
JEL: D82; D83; C72

*We thank the Neubauer Family Foundation, the William Ladany Faculty Research Fund, and the Initiative on Global Markets at the University of Chicago Booth School of Business for financial support.

1 Introduction

In many economic settings, informed experts choose how much of their private information to disclose to a decision-maker (DM) who will take an action that affects the payoffs of all the players. Often, the disclosed information is certifiable, meaning that the claims made by the experts might be more or less informative but cannot be false.¹ A large literature establishes conditions under which all private information will be disclosed in equilibrium.² A key insight from this literature is that when experts' preferences are suitably monotonic in DM's beliefs or sufficiently opposed, full disclosure is an equilibrium. With a single expert who has monotonic preferences, this equilibrium is unique.

Models in this literature typically take the experts' initial private information as exogenous. This is a suitable assumption in many settings. There are other settings, however, where information is symmetric at the outset and the experts choose how much private information to gather. For example, a pharmaceutical company may or may not conduct clinical trials that assess a drug's efficacy for a particular demographic group. Whenever private information is costless and can be covertly gathered, it is a dominant strategy to become as informed as possible. When the acquisition of private information is public, however, becoming more informed may be harmful in equilibrium. If the FDA knows that a pharmaceutical company conducted a clinical trial specifically to assess a drug's side effects in children, the failure to disclose the results is likely to generate skepticism.

In this paper, we study disclosure when private information is endogenous. We consider *ex ante* symmetric information games where $n \geq 1$ experts simultaneously conduct experiments about the state of the world. More informative experiments are (weakly) more costly to the expert. DM observes which experiments are conducted, and each expert privately observes the outcome of his own experiment. The experts convey verifiable messages to DM about the outcomes. DM then takes an action. We focus on pure-strategy perfect Bayesian equilibria. An *outcome* of the game is the joint distribution of the state of the world, DM's beliefs, DM's actions, and all the players' payoffs.

We study the effect of requiring experts to fully disclose the results of their experiments. We might expect such regulations to change the equilibrium outcomes when the usual monotonicity or opposed preferences conditions for full disclosure are not satisfied. This could benefit the decision

¹In the persuasion games literature, terms *certifiable*, *verifiable*, and *provable* are typically used interchangeably. A separate literature initiated by Crawford and Sobel (1982) examines cheap talk communication.

²For seminal contributions, see Grossman (1981), Milgrom (1981), and Milgrom & Roberts (1986). For a recent survey, see Milgrom (2008).

maker if it causes more information to be revealed. It could also benefit the experts if the inability to commit to truthful disclosure reduces their equilibrium payoffs, as can happen in cheap talk settings.

Our main result is that disclosure requirements have no effect on the set of equilibrium outcomes and thus no effect on either the decision-maker's or the experts' payoffs. Essentially, we show that endogenous information will always be disclosed in equilibrium; if there is an equilibrium in which information is withheld, the outcome must be the same as in another equilibrium with truthful disclosure. Moreover, if strictly more informative signals are strictly more costly, there is no equilibrium where information is withheld.

The remainder of the paper is structured as follows. Section 2 covers some mathematical preliminaries. Section 3 presents the model. The statement and the proof of the main result are in section 4. We discuss the relationship to the existing literature in the final section.

2 Mathematical preliminaries

In this section we introduce notation and mathematical concepts that are building blocks of our model. Both the notation and the particular way of formalizing signals are taken from Gentzkow and Kamenica (2012).

2.1 State space and signals

Let Ω be a finite state space. A state of the world is denoted by $\omega \in \Omega$. A belief is denoted by μ . The prior distribution on Ω is denoted by μ_0 . Let X be a random variable that is independent of ω and uniformly distributed on $[0, 1]$ with typical realization x . We model signals as deterministic functions of ω and x . Formally, a *signal* π is a finite partition of $\Omega \times [0, 1]$ s.t. $\pi \subset S$, where S is the set of non-empty Lebesgue measurable subsets of $\Omega \times [0, 1]$. We refer to any element $s \in S$ as a *signal realization*.

With each signal π we associate an S -valued random variable that takes value $s \in \pi$ when $(\omega, x) \in s$. Let $p(s|\omega) = \lambda(\{x | (\omega, x) \in s\})$ and $p(s) = \sum_{\omega \in \Omega} p(s|\omega) \mu_0(\omega)$ where $\lambda(\cdot)$ denotes the Lebesgue measure. For any $s \in \pi$, $p(s|\omega)$ is the conditional probability of s given ω and $p(s)$ is the unconditional probability of s .

2.2 Lattice structure

The formulation of a signal as a partition induces an algebraic structure on the set of signals. This structure allows us to “add” signals together and thus easily examine their joint information content. Let Π be the set of all signals. Let \supseteq denote the refinement order on Π , that is, $\pi_1 \supseteq \pi_2$ if every element of π_1 is a subset of an element of π_2 . The pair (Π, \supseteq) is a lattice. The join $\pi_1 \vee \pi_2$ of two elements of Π is defined as the supremum of $\{\pi_1, \pi_2\}$. For any finite set (or vector) P we denote the join of all its elements by $\vee P$. We write $\pi \vee P$ for $\pi \vee (\vee P)$. Note that $\pi_1 \vee \pi_2$ is a signal that consists of signal realizations s such that $s = s_1 \cap s_2$ for some $s_1 \in \pi_1$ and $s_2 \in \pi_2$. Hence, $\pi_1 \vee \pi_2$ is the signal that yields the same information as observing both signal π_1 and signal π_2 . In this sense, the binary operation \vee “adds” signals together.

2.3 Distributions of posteriors

A *distribution of posteriors*, denoted by τ , is an element of $\Delta(\Delta(\Omega))$ that has finite support.³ Observing a signal realization s s.t. $p(s) > 0$ generates a unique posterior belief⁴

$$\mu_s(\omega) = \frac{p(s|\omega)\mu_0(\omega)}{p(s)} \text{ for all } \omega.$$

We write $\langle \pi \rangle$ for the distribution of posteriors induced by π . It is easy to see that $\tau = \langle \pi \rangle$ assigns probability $\tau(\mu) = \sum_{\{s \in \pi: \mu_s = \mu\}} p(s)$ to each μ .

3 The model

3.1 The baseline game

The decision-maker (DM) has a continuous utility function $u(a, \omega)$ that depends on her action $a \in A$ and the state of the world $\omega \in \Omega$. There are $n \geq 1$ experts indexed by i . Each expert i has a continuous utility function $v_i(a, \omega)$ that depends on DM’s action and the state of the world. All experts and DM share the prior μ_0 . The action space A is compact.

The timing is as follows:

³The fact that distributions of posteriors have finite support follows from the assumption that each signal has finitely many realizations. The focus on such signals is without loss of generality under the maintained assumption that Ω is finite.

⁴For those s with $p(s) = 0$, set μ_s to be an arbitrary belief.

1. Each expert simultaneously chooses a signal π_i , the choice of which is not observed by the other experts.
2. Each expert privately observes the realization s_i of his own signal.
3. Each expert simultaneously sends a verifiable message $m_i \in M(s_i)$ to DM.
4. DM observes the signals chosen by the experts and all of the messages.
5. DM chooses an action.

The set $M(s)$ specifies the set of messages that are feasible upon observing signal realization s . Let $M = \cup_{s \in S} M(s)$ denote the set of all possible messages. For each $m \in M$, let $T(m) = \{s \in S | m \in M(s)\}$. We say that message m *verifies* s if $T(m) = \{s\}$. We assume that for each $s \in S$ there exists a unique message that verifies it.⁵

For each expert i , we let $c_i : \Pi \rightarrow \bar{\mathbb{R}}_+$ denote the cost of each signal.⁶ Expert i 's payoff in state ω is thus $v_i(a, \omega) - c_i(\pi)$ if he chooses signal π and the decision-maker takes action a . We assume that more informative signals are more expensive: $\pi \succeq \pi' \Rightarrow c_i(\pi) \geq c_i(\pi')$ for any $\pi, \pi' \in \Pi$ and any i . This is an important assumption that is absolutely necessary for our result. If an expert can save money by becoming more informed and then withholding the additional information, full disclosure will not happen in equilibrium and thus disclosure requirements will change the set of equilibrium outcomes.

Let $\sigma_i = (\pi_i, (\gamma_i^\pi)_{\pi \in \Pi})$ denote expert i 's strategy, and σ denote a strategy profile. A strategy for expert i consists of a signal $\pi_i \in \Pi$ and a feasible messaging policy⁷ $\gamma_i^\pi : S \rightarrow \Delta(M)$ following each signal π .⁸ We say σ_i is a pure-strategy if the messaging policy is deterministic on the equilibrium path (i.e., $\gamma_i^{\pi_i}$ is deterministic). Let $\tilde{\mu}(\mathbf{m}) \in \Delta(S^n)$ denote DM's belief, given the vector of messages \mathbf{m} , about the signal realizations observed by the experts. The structure of the information sets requires DM's belief about expert i 's signal to have support in $T(m_i)$ (the set of i 's "types" for which m_i was an available message). Since DM knows the experts' choices of signals, each belief $\tilde{\mu}(\mathbf{m})$ about the signal realizations implies a unique belief μ about ω . Throughout the paper we assume that DM has a unique optimal action at any given belief, i.e., $a^*(\mu) \equiv \operatorname{argmax}_{a \in A} E_\mu[u(a, \omega)]$ is

⁵The uniqueness assumption is not needed for our result. It simplifies our notation, however, by making the truthful messaging policy unique.

⁶ $\bar{\mathbb{R}}_+$ denotes the affinely extended non-negative real numbers: $\bar{\mathbb{R}}_+ = \mathbb{R} \cup \{\infty\}$. Allowing the cost to be infinite for some signals incorporates the cases where not all the experts have access to all the signals.

⁷A messaging policy γ_i^π is feasible if $\operatorname{supp}(\gamma_i^\pi(s)) \subset M(s)$ for all s .

⁸As we focus on pure-strategy equilibria, we do not introduce notation for mixed strategies in the choice of π_i .

single-valued for all μ . By the theorem of the maximum, the fact that $a^*(\cdot)$ is single-valued implies that it is continuous.

Let $\mathcal{B}^i(\boldsymbol{\sigma}_{-i}, \tilde{\mu})$ denote the best-response correspondence for expert i . A strategy σ_i is in this set if and only if it is a best response to $\boldsymbol{\sigma}_{-i}$ and $\tilde{\mu}$ at every information set. Expert i 's information sets are the ‘‘initial’’ node where he selects a signal, and each possible (π, s) s.t. $\pi \in \Pi$ and $s \in \pi$. Note that this best-response correspondence does not depend on DM's strategy; because DM has a unique optimal action at every $\tilde{\mu}(\mathbf{m})$, expert i can take her behavior (given $\tilde{\mu}$) as fixed. A pair $(\boldsymbol{\sigma}, \tilde{\mu})$ is a (perfect Bayesian) equilibrium if and only if $\tilde{\mu}$ obeys Bayes' rule on the equilibrium path and $\sigma_i \in \mathcal{B}^i(\boldsymbol{\sigma}_{-i}, \tilde{\mu})$ for all i . We define the *outcome* of the game to be the joint distribution of the state of the world, DM's beliefs, DM's actions, and all the players' payoffs.

A pure-strategy σ_i defines a partition π' of $\Omega \times [0, 1]$ with each m_i sent in equilibrium corresponding to one element of the partition. Thus, if DM knows expert i is playing σ_i , the distribution of posteriors induced by her observation of his messages is $\langle \pi' \rangle$. We denote this signal equivalent to σ_i by $r(\sigma_i)$. Note that $\pi_i \supseteq r(\sigma_i)$, which implies that $c_i(\pi_i) \geq c_i(r(\sigma_i))$. Given $\boldsymbol{\sigma} = (\sigma_1, \dots, \sigma_n)$, let $\mathbf{r}(\boldsymbol{\sigma})$ denote $(r(\sigma_1), \dots, r(\sigma_n))$. Given an equilibrium $(\boldsymbol{\sigma}, \tilde{\mu})$, its outcome is $\langle \vee \mathbf{r}(\boldsymbol{\sigma}) \rangle$.

3.2 Disclosure requirements

We define an alternative game with *disclosure requirements* in which experts are required to disclose their private information truthfully. To distinguish this from the game defined above, we refer to the latter as the *baseline game*.

Let γ^* denote the truthful messaging policy, i.e., $\gamma^*(s)$ places probability 1 on the message that verifies s . Under disclosure requirements, each expert i must set $\gamma_\pi = \gamma^*$ for every $\pi \in \Pi$. Accordingly, we can represent each expert's strategy simply as π_i and let $\boldsymbol{\pi}$ denote a strategy profile.

Let $\mathcal{B}_{DR}^i(\boldsymbol{\pi}_{-i})$ denote the best-response correspondence for expert i under disclosure requirements. A strategy profile $\boldsymbol{\pi}$ is a pure-strategy equilibrium under disclosure requirements if and only if $\pi_i \in \mathcal{B}_{DR}^i(\boldsymbol{\pi}_{-i})$. Given an equilibrium $\boldsymbol{\pi}$, its *outcome* is $\langle \vee \boldsymbol{\pi} \rangle$.

4 Main result

Our main result is the following:

Theorem 1. *Disclosure requirements do not change the set of pure-strategy equilibrium outcomes.*

The statement of this result does not require any assumptions about the number of experts nor about the experts' or the decision-maker's utility functions. Moreover, the theorem does not only state that even without disclosure requirements there exists an equilibrium with full revelation of private information; rather it establishes a stronger claim that disclosure requirements have no impact whatsoever on the entire set of equilibrium outcomes.

The remainder of this section provides a proof of Theorem 1. Let Σ^π denote the set of strategies that select signal π . Let Σ^* denote the set of strategies that utilize truthful messaging on the equilibrium path, i.e., strategies of the form $(\pi', (\gamma^\pi)_{\pi \in \Pi})$ s.t. $\gamma^{\pi'} = \gamma^*$. To show that any outcome under disclosure requirements is also an outcome of the baseline game, it will suffice to establish the following Lemma.

Lemma 1. *Given some π , suppose $\pi_i \in \mathcal{B}_{DR}^i(\pi_{-i})$ for all i . Then, there exist σ and $\tilde{\mu}$ such that*

- (i) $\tilde{\mu}$ obeys Bayes' rule given σ
- (ii) $\sigma_i \in \mathcal{B}^i(\sigma_{-i}, \tilde{\mu})$ for all i
- (iii) $\sigma_i \in \Sigma^* \cap \Sigma^{\pi_i}$ for all i

Proof. We begin the proof by introducing an auxiliary game $G^i(\pi_S, \pi_R)$. There is a single expert with utility v_i and DM with utility u . The timing is: (i) DM privately observes a signal realization s_R from signal π_R ; (ii) the expert privately observes a signal realization s_S from signal π_S ; (iii) the expert sends a message $m \in M(s_S)$ to DM; (iv) DM takes an action. Let γ denote the expert's messaging strategy and $\eta(m)$ denote DM's beliefs, given m , about s_S . An equilibrium of this game is a pair (γ, η) s.t. η obeys Bayes' rule on the equilibrium path and, at each information set s_S , γ_i is the best response to η . (Since DM has a unique optimal action at every μ , and thus at every $\eta(m)$, we do not need to specify her behavior given her belief). Standard arguments ensure that, given any π_S and π_R , the game $G^i(\pi_S, \pi_R)$ has a perfect Bayesian equilibrium.

To construct the requisite σ and $\tilde{\mu}$, we begin by imposing condition (iii) in the Lemma, i.e., for each expert i we specify that σ_i selects π_i at the initial information set and that $\gamma^{\pi_i} = \gamma^*$. We also begin the construction of $\tilde{\mu}$ by imposing condition (i): $\tilde{\mu}$ follows Bayes rule given π followed by truthful messaging by all of the experts.

We next construct off-equilibrium messaging policies $(\gamma_i^\pi)_{\pi \neq \pi_i}$ for each expert. Consider expert i . Given any $\pi \in \Pi$, consider the auxiliary game $G^i(\pi, \bigvee_{j \neq i} \pi_j)$, which has some equilibrium (γ', η) . We set γ^π to γ' and we set $\tilde{\mu}((\pi, \pi_{-i}), (m_i, \mathbf{m}_{-i}))$ for equilibrium \mathbf{m}_{-i} as follows: DM's belief about signal realizations observed by other experts is already pinned down (since other experts are

playing truthful messaging policies), and we set DM's belief about the signal realization observed by expert i to $\eta(m_i)$. We repeat this procedure for each $\pi \neq \pi_i$. This completes the construction of σ_i and specifies $\tilde{\mu}$ on all DM's information sets off the equilibrium path where only expert i deviates from σ_i . We can then repeat this procedure expert by expert and thus construct the entire strategy profile σ , as well as $\tilde{\mu}$ on all DM's information sets where only one expert deviates from σ . We can choose an arbitrary specification of $\tilde{\mu}$ on DM's information sets where multiple experts deviate from σ .

These σ and $\tilde{\mu}$ satisfy conditions (i) and (iii) of the Lemma by construction. For any expert i , any $\pi \neq \pi_i$, and any $s \in \pi$, condition (ii) is satisfied on information set (π, s) because γ_i^π is an equilibrium messaging policy of $G^i(\pi, \bigvee_{j \neq i} \pi_j)$. To show that condition (ii) is satisfied on the equilibrium path (on information sets (π_i, s)), we need to establish that γ^* is an equilibrium messaging policy of $G^i(\pi_i, \bigvee_{j \neq i} \pi_j)$. We know $G^i(\pi_i, \bigvee_{j \neq i} \pi_j)$ has some equilibrium, say (γ, η) . It will suffice to show that given η , expert i 's payoff from γ^* is the same as his payoff from γ following any s . Denote these payoffs by $y^*(s)$ and $y(s)$, respectively. Since (γ, η) is an equilibrium, we know $y(s) \geq y^*(s) \forall s \in \pi_i$. Moreover, since $\pi_i \in \mathcal{B}_{DR}^i(\pi_{-i})$, we know $\sum_{\omega \in \Omega} \sum_{s \in \pi_i} y^*(s) p(s|\omega) \mu_0(\omega) \geq \sum_{\omega \in \Omega} \sum_{s \in \pi_i} y(s) p(s|\omega) \mu_0(\omega)$. (Otherwise, under disclosure requirements expert i could profitably deviate from π_i to a signal which garbles π_i by γ .) Combining these two inequalities yields $y(s) = y^*(s)$.

It remains to show that condition (ii) is satisfied for each expert at the initial information set where he chooses his signal. Let $\hat{v}_i(\mu) \equiv \mathbb{E}_\mu v_i(a^*(\mu), \omega)$. Let v_i^* be expert i 's payoff under σ and $\tilde{\mu}$. Since $\pi_i \in \mathcal{B}_{DR}^i(\pi_{-i})$, we know

$$v_i^* \geq \mathbb{E}_{\langle \pi \vee \pi_{-i} \rangle} [\hat{v}_i(\mu)] - c_i(\pi) \quad \forall \pi \in \Pi. \quad (1)$$

Suppose expert i deviates from $\sigma_i = (\pi_i, (\gamma_i^\pi)_\pi)$ to $\sigma'_i = (\pi'_i, (\gamma_i^\pi)_\pi)$. By our construction of $\tilde{\mu}$ through η , we know the distribution of DM's posterior must be $\langle r(\sigma'_i) \vee \pi_{-i} \rangle$. Hence, the deviation yields the payoff

$$\mathbb{E}_{\langle r(\sigma'_i) \vee \pi_{-i} \rangle} [\hat{v}_i(\mu)] - c_i(\pi'_i) \leq \mathbb{E}_{\langle r(\sigma'_i) \vee \pi_{-i} \rangle} [\hat{v}_i(\mu)] - c_i(r(\sigma'_i)) \leq v_i^*$$

where the first inequality follows from the fact that $\pi'_i \supseteq r(\sigma'_i)$ implies $c_i(\pi'_i) \geq c_i(r(\sigma'_i))$ and the second inequality follows from Equation (1). Since the deviation yields a weakly lower payoff than v_i^* , condition (ii) is also satisfied at the initial information set. \square

To show that any outcome of the baseline game is also an outcome under disclosure requirements, we begin with the following Lemma.

Lemma 2. *Suppose $(\sigma, \tilde{\mu})$ is an equilibrium. Then, $r(\sigma_i) \in \mathcal{B}_{DR}^i(\mathbf{r}(\sigma_{-i}))$ for all i .*

Proof. Consider any expert i . His equilibrium strategy is some $\sigma_i = (\pi_i, (\gamma^\pi)_\pi)$. Let v^* denote his equilibrium payoff. Given $(\sigma_{-i}, \tilde{\mu})$, for every signal $\pi' \in \Pi$, let $v_{\pi'}$ denote his payoff if he deviates to strategy $(\pi', (\gamma^\pi)_\pi)$, and let $v_{\pi'}^*$ denote his payoff if he deviates to strategy $(\pi', (\gamma^*)_\pi)$. Since $(\sigma, \tilde{\mu})$ is an equilibrium, we know: (i) $v^* \geq v_{\pi'}$ for all π' (by the fact that σ_i was the best response at the initial information set); and (ii) $v_{\pi'} \geq v_{\pi'}^*$ for all π' (by the fact that σ_i was the best response at each information set (π', s)). Finally, if expert i deviates to strategy $(r(\sigma_i), (\gamma^*)_\pi)$, his payoff under this deviation must also be v^* , so $v_{r(\sigma_i)}^* = v^*$. Combining this with inequalities (i) and (ii) we obtain $v_{r(\sigma_i)}^* \geq v_{\pi'}^*$ for all π' . Since for all π' , $v_{\pi'}^* = \mathbb{E}_{\langle \pi' \vee \mathbf{r}(\sigma_{-i}) \rangle} [\hat{v}_i(\mu)]$, this implies $r(\sigma_i) \in \mathcal{B}_{DR}^i(\mathbf{r}(\sigma_{-i}))$. \square

This Lemma shows that, given any equilibrium of the baseline game, there is an equilibrium under disclosure requirements that induces the same joint distribution of the state of the world, DM's beliefs, DM's actions, and DM's payoffs. It only remains to add experts' costs of signals to this list. The only way that these costs could be different is some expert i were utilizing a strategy $\sigma_i = (\pi_i, (\gamma_i^\pi)_{\pi \in \Pi})$ s.t. $c_i(\pi_i) > c_i(r(\sigma_i))$. But this could not happen in equilibrium since it would then be profitable for expert i to deviate to a strategy in $\Sigma^{r(\sigma_i)} \cap \Sigma^*$.

This completes the proof of Theorem 1.

5 Related Literature

5.1 Persuasion games with verifiable types

As mentioned in the introduction, a large literature examines disclosure of exogenous private information in persuasion games, i.e., settings where informed expert(s) can send verifiable messages. Milgrom (1981) shows that full disclosure is a unique equilibrium outcome when there is a single expert who can send any verifiable message and whose preferences are monotonic (whether the expert, who knows the true state is ω^* , prefers DM to believe the state is ω or ω' does not depend on ω^*). Since this early contribution, this literature has evolved along three distinct dimensions.

Weakening monotonicity. Seidmann and Winter (1997), Giovannoni and Seidmann (2007), and Mathis (2008) establish that existence of a fully revealing equilibrium can be guaranteed if

we replace the monotonicity assumption with a somewhat weaker single-crossing property (if the expert, when he knows the true state is ω^* , prefers DM to believe the state is ω rather than $\omega' \leq \omega$, the expert also has this preference when he knows the true state is $\omega^{**} > \omega^*$). Moreover, if (in addition) the preference conflict is “stable” (e.g., at any ω^* , expert’s ideal action by DM is always greater than DM’s ideal action), then the fully revealing equilibrium is unique. Hagenbach *et al.* (2012) introduce a general model that encompasses much of this literature and establish a simple condition that is necessary and sufficient for existence of a fully revealing equilibrium.

Weakening verifiability. Milgrom (1981) assumes that set of messages is the power set of the experts’ type. Okuno-Fujiwara *et al.* (1990) and others point out that other message spaces can be assumed and that full revelation can remain the unique outcome even if the expert cannot always verify his type. For example, it would not matter if the expert could not prove that he is a “low” type. In spirit of these results, we put limited structure on sets $M(s)$ and only impose the key assumption that for each s there is a message that verifies it.

Introducing multiple experts. In all of the aforementioned papers, full revelation is driven by some version of the “unraveling argument” – if some types pool, at least one of them is “better” than the “average” and will prefer to reveal himself. When there are multiple experts, however, there are other forces that can lead to full revelation.⁹ If for each state there is some expert who wishes to disclose the state to DM so as to avoid her default action, full revelation is an equilibrium (Milgrom and Roberts 1986). Also, Lipman and Seppi (1995) establish that, as long as DM knows experts have conflicting preferences, there is a full revelation equilibrium, even under limited verifiability.

Under our assumption of endogenous information (and complete verifiability), we show that for any number of experts and for any configuration of preferences (regardless of monotonicity, single-crossing, or conflict) full revelation of private information is always an equilibrium. Moreover, given any equilibrium, there is a full revelation equilibrium that induces the same outcomes.

In most models of verifiable communication, there is a separate set of experts who wish to influence a third party (DM). That said, some papers examine environments where experts disclose private information and then play games with each other (Okuno-Fujiwara *et al.* 1990, Hagenbach and Koessler 2011, Hagenbach *et al.* 2012). When private information is exogenous, this distinction is not particularly important – it does not matter whether the publicly disclosed information impacts experts’ payoff through an action of a third-party DM or through the equilibrium outcome of the

⁹Okuno-Fujiwara *et al.* (1990) and Hagenbach and Koessler (2011) apply unraveling-like arguments to settings with multiple experts.

post-disclosure game. Our results, however, only apply to the environments where experts seek to influence a third party. Once experts' information is endogenous, the publicly disclosed information is no longer sufficient to determine the payoffs of a post-disclosure game.

Finally, existing literature considers both settings where experts are informed about a common state of the world (e.g., Milgrom and Roberts 1986, Lipman and Seppi 1995) and settings where each expert has private information only about his own type (Okuno-Fujiwara *et al.* 1990, Hagenbach and Koessler 2011, Hagenbach *et al.* 2012). Our model covers both of these case. If we let $c_i(\cdot)$ be the same for all experts (e.g., $c_i(\pi) = 0$ for all π), then all experts can become privately informed about the “common” state ω . Alternatively, suppose that $\Omega = T_1 \times \dots \times T_n$ where T_i is the set of possible types of expert i . Then, we can set $c_i(\pi) = \infty$ for any π that is informative about types other than i and thus capture the setting where an expert can only become privately informed about his own type – he cannot learn about nor disclose to DM any information about any other expert.

5.2 Competition in persuasion

Kamenica and Gentzkow (2011) analyze a game where a single expert wishes to influence DM's action by choosing an observable costless signal about the state the world. They focus on identifying the conditions under which the expert benefits from the ability to generate such a signal, and on characterizing the distribution of DM's beliefs under the optimal signal. In the discussion of their model, they point out that the outcome of their game would be the same if DM did not observe the signal realizations directly but the expert could send verifiable messages. This observation is the starting point of our analysis.

Gentzkow and Kamenica (2012) examine a game where any number of experts choose costless signals about a state of the world, DM directly observes the realizations of the signals, and then DM takes an action that affects the welfare of all the players. This game is identical to the game faced by the experts in our model under disclosure regulation.¹⁰ Theorem 1 thus implies that the set of equilibrium outcomes of the game we consider, where experts convey verifiable messages and can withhold unfavorable information *ex post*, coincides with the set of equilibrium outcomes of the

¹⁰Gentzkow and Kamenica (2012) assume that each expert i can select a signal whose realization (conditional on ω) is arbitrarily correlated with the realizations of other experts' signals. This is the case in our model if we set $c_i(\pi) = 0$ for all $\pi \in \Pi$. Our model also encompasses environments, however, where such correlation is not possible. Specifically, let $Y = Y_1 \times \dots \times Y_n$ where $Y_i = [0, 1]$ for all i and let f be a bijection from Y to $[0, 1]$. Then, we could set $c_i(\pi) = \infty$ if there exist $s, s' \in \pi$ s.t. $s \cap f((y_i, y_{-i})) \neq s' \cap f((y_i, y'_{-i}))$ for some $(y_i, y_{-i}), (y_i, y'_{-i}) \in Y$. In other words, we could redefine X as an n -dimensional random variable and we let the signal realization of expert i only depend on the i th dimension of X .

observable signal game considered by Gentzkow and Kamenica (2012).¹¹ Their paper characterizes the set of equilibrium outcome and derives comparative statics of the outcomes' informativeness with respect to the extent of competition. Since their results are all about the set of equilibrium outcomes (rather than equilibrium strategies), Theorem 1 means that their results apply to our model as well.

¹¹Our proof of Theorem 1 requires that DM's optimal action be unique at every belief, an assumption not imposed by Gentzkow and Kamenica (2012). When this assumption is not satisfied, the equivalence of the two games can be guaranteed by introducing a small amount of private information for DM, so that the distribution of DM's optimal actions is single-valued and continuous.

References

- Crawford, Vincent P., and Joel Sobel. 1982. Strategic information transmission. *Econometrica*. 50:1431-51.
- Gentzkow, Matthew, and Emir Kamenica. 2012. Competition in persuasion. Working paper.
- Giovannoni, Francesco, and Daniel J. Seidmann. 2007. Secrecy, two-sided bias and the value of evidence. *Games and economic behavior*. 59:296-315.
- Grossman, Sanford J. 1981. The informational role of warranties and private disclosure about product quality. *Journal of Law and Economics*. 24:461-83.
- Hagenbach, Jeanne, and Frederic Koessler. 2011. Full disclosure in decentralized organizations. Working paper.
- Hagenbach, Jeanne, Frederic Koessler, and Eduardo Perez-Richet. 2012. Certifiable pre-play communication: Full disclosure. Working paper.
- Kamenica, Emir, and Matthew Gentzkow. 2011. Bayesian persuasion. *American Economic Review*. 101:2590-615.
- Lipman, Barton L., and Duane J. Seppi. 1995. Robust inference in communication games with partial provability. *Journal of Economic Theory*. 66:370-405.
- Mathis, Jerome. 2008. Full revelation of information in Sender-Receiver games of persuasion. *Journal of Economic Theory*. 143:571-84.
- Milgrom, Paul. 1981. Good news and bad news: Representation theorem and applications. *The Bell Journal of Economics*. 12:380-91.
- Milgrom, Paul. 2008. What the seller won't tell you: Persuasion and disclosure in markets. *Journal of Economic Perspectives*. 22:115-31.
- Milgrom, Paul, and John Roberts. 1986. Relying on the information of interested parties. *Rand Journal of Economics*. 17:18-32.
- Okuno-Fujiwara, Masahiro, Andrew Postlewaite, and Kotaro Suzumura. 1990. Strategic information revelation. *Review of Economic Studies*. 57:25-47.
- Seidmann, Daniel J., and Eyal Winter. 1997. Strategic information transmission with verifiable messages. *Econometrica*. 65:163-9.